

# The Solovay-Kitaev Theorem

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# 1 Preliminaries

## 1.1 Spectral Norm

We define the spectral norm in  $SU(2)$  by

$$\|A\| = \sup_{v \in \mathbb{C}^2 \setminus \{0\}} \frac{\|Av\|}{\|v\|}$$

using the natural norm on  $\mathbb{C}$ . We note that equivalently, we obtain

$$\|A\| = \sigma_{\max} = \sqrt{\lambda_{\max}(A^\dagger A)}$$

where  $\sigma_{\max}$  is the largest singular value of  $A$ . We also note that for a normal matrix  $N$ ,

$$\|N\| = |\lambda_{\max}(N)|.$$

## 1.2 Universal Gate Sets

We say that a finite subset  $\Gamma \subseteq SU(2)$  is a universal gate set if  $\langle \Gamma \rangle$  is dense in  $SU(2)$  and  $\Gamma$  is closed under inverses.

# 2 Shrinking Lemma

A key component of the proof is to be able to reduce the size of our net around the identity of  $SU(2)$ . By repeatedly taking the commutator of pairs in our original net, we find that we can construct arbitrarily close nets and hence approximate elements of  $SU(2)$ .

We note that any the Pauli matrices  $\sigma_x, \sigma_y, \sigma_z$  form a basis for  $\mathfrak{su}(2)$ . Hence we can represent any element  $A \in SU(2)$  by  $A = e^{i\vec{a} \cdot \vec{\sigma}}$  where  $\vec{a} \in \mathbb{R}^3$  and  $\vec{\sigma}$  is a 3-vector of Pauli matrices. Without loss of generality, we can take  $\|\vec{a}\| \leq \pi$ .

## 2.1 Facts about $SU(2)$

Repeatedly, we will find ourselves reducing distances in  $SU(2)$  to those in  $\mathbb{R}^3$ . This allows us to make use of concepts such as the cross-product and trigonometric functions. Hence we provide some facts define the relation between the two metrics.

**Fact 1.** If  $\vec{a} \in \mathbb{R}^3$  then  $\|e^{i\vec{a} \cdot \vec{\sigma}} - I\| = 2 \sin\left(\frac{\|\vec{a}\|}{2}\right) = \|\vec{a}\| + O(\|\vec{a}\|^3)$ .

*Proof.* Pick  $\vec{a} \in \mathbb{R}^3$ . We first wish to compute the eigenvalues of

$$\vec{a} \cdot \vec{\sigma} = \begin{bmatrix} a_z & a_x - ia_y \\ a_x + ia_y & -a_z \end{bmatrix}$$

which we can do in the traditional way.

$$\begin{aligned}
\begin{vmatrix} a_z - \lambda & a_x - ia_y \\ a_x + ia_y & -a_z - \lambda \end{vmatrix} &= -(a_z - \lambda)(a_z + \lambda) - (a_x - ia_y)(a_x + ia_y) \\
&= -(a_z^2 - \lambda^2) - (a_x^2 - a_y^2) \\
&= \lambda^2 - (a_x^2 + a_y^2 + a_z^2) \\
&= \lambda^2 - \|\vec{a}\|^2.
\end{aligned}$$

Hence our eigenvalues are  $\lambda = \pm \|\vec{a}\|$ .

We now observe that any eigenvalue of  $\vec{a} \cdot \vec{\sigma}$  is an eigenvalue of  $e^{i\vec{a} \cdot \vec{\sigma}}$  with the same eigenvector which follows from the power series definition of matrix exponents. As  $e^{i\vec{a} \cdot \vec{\sigma}}$  is also a  $2 \times 2$  matrix, we have that all the eigenvalues of  $e^{i\vec{a} \cdot \vec{\sigma}}$  are of the form  $e^{i\lambda}$  for eigenvalues  $\lambda$  of  $\vec{a} \cdot \vec{\sigma}$ . Thus the eigenvalues of  $e^{i\vec{a} \cdot \vec{\sigma}}$  are  $e^{\pm i\|\vec{a}\|}$ .

It then follows that the eigenvalues of  $e^{i\vec{a} \cdot \vec{\sigma}} - I$  are  $e^{\pm i\|\vec{a}\|} - 1$ . We also know that because  $e^{i\vec{a} \cdot \vec{\sigma}}$  and  $I$  commute and are both normal, that their sum is normal. Hence we have that the singular values are the absolute value of the eigenvalues. We then compute the singular values.

$$\begin{aligned}
|e^{\pm i\|\vec{a}\|} - 1| &= \sqrt{(\cos(\|\vec{a}\|) - 1)^2 + (\sin(\|\vec{a}\|))^2} \\
&= \sqrt{2 - 2\cos(\|\vec{a}\|)} \\
&= \sqrt{4\sin^2(\|\vec{a}\|/2)} \\
&= 2\sin\left(\frac{\|\vec{a}\|}{2}\right).
\end{aligned}$$

Thus  $\|e^{i\vec{a} \cdot \vec{\sigma}} - I\| = 2\sin\left(\frac{\|\vec{a}\|}{2}\right)$ . We proceed with the power series expansion of  $\sin(x)$ .

$$\begin{aligned}
2\sin\left(\frac{\|\vec{a}\|}{2}\right) &= 2\left(\frac{\|\vec{a}\|}{2} - \frac{(\|\vec{a}\|/2)^3}{3!} + \frac{(\|\vec{a}\|/2)^5}{5!} - \dots\right) \\
&= \|\vec{a}\| - \frac{(2\|\vec{a}\|/2)^3}{3!} + \frac{(2\|\vec{a}\|/2)^5}{5!} - \dots \\
&= \|\vec{a}\| + O(\|\vec{a}\|^3).
\end{aligned}$$

□

**Fact 2.** If  $\vec{b}, \vec{c} \in \mathbb{R}$  then  $\|e^{i\vec{b} \cdot \vec{\sigma}} - e^{i\vec{c} \cdot \vec{\sigma}}\| \leq \|\vec{b} - \vec{c}\|$ .

*Proof.* By applying [Fact 1](#) we obtain

$$\|e^{i\vec{b} \cdot \vec{\sigma}} - e^{i\vec{c} \cdot \vec{\sigma}}\| = \|e^{i(\vec{b}-\vec{c}) \cdot \vec{\sigma}} - I\| = 2\sin\left(\frac{\|\vec{b} - \vec{c}\|}{2}\right) \leq \|\vec{b} - \vec{c}\|.$$

□

**Fact 3.** If  $\vec{b}, \vec{c} \in \mathbb{R}$  then  $\left[ \vec{b} \cdot \vec{\sigma}, \vec{c} \cdot \vec{\sigma} \right] = 2i(\vec{b} \times \vec{c}) \cdot \vec{\sigma}$ .

*Proof.* We note the following properties of Pauli matrices and the cross product.

$$\begin{aligned} [\sigma_j, \sigma_k] &= 2i\varepsilon_{jkl}\sigma_l, \\ (e_j \times e_k) &= \varepsilon_{jkl}e_l. \end{aligned}$$

Consider an arbitrary component of our commutator:

$$\begin{aligned} \left[ \vec{b} \cdot \vec{\sigma}, \vec{c} \cdot \vec{\sigma} \right]_l &= b_j c_k [\sigma_j, \sigma_k] + b_k c_j [\sigma_k, \sigma_j] \\ &= b_j c_k [\sigma_j, \sigma_k] - b_k c_j [\sigma_j, \sigma_k] \\ &= (b_j c_k - b_k c_j) [\sigma_k, \sigma_j] \\ &= 2i(b_j c_k - b_k c_j) \varepsilon_{jkl} \sigma_l \\ &= 2i(\vec{b} \times \vec{c})_l \sigma_l. \end{aligned}$$

Thus we have our result. □

**Fact 4.** If  $\vec{b}, \vec{c} \in \mathbb{R}$  with  $\|\vec{b}\| = O(\varepsilon)$  and  $\|\vec{c}\| = O(\varepsilon)$  then  $\left\| \left[ e^{i\vec{b} \cdot \vec{\sigma}}, e^{i\vec{c} \cdot \vec{\sigma}} \right] - e^{-[\vec{b} \cdot \vec{\sigma}, \vec{c} \cdot \vec{\sigma}]} \right\| = O(\varepsilon^3)$ .

*Proof.* Follows from Ozols' proof [1]. □

**Lemma 1.** (*Shrinking Lemma*) If  $\Gamma$  is an  $\varepsilon^2$ -net for  $S_\varepsilon$  then  $[\Gamma, \Gamma]$  is an  $O(\varepsilon^3)$ -net for  $S_{\varepsilon^2}$ .

*Proof.* Suppose  $\Gamma$  is an  $\varepsilon^2$  net for  $S_\varepsilon$ . We want to show for any  $A \in S_{\varepsilon^2}$ , there exist  $U, V \in \Gamma$  such that  $\|A - [U, V]\| = O(\varepsilon^3)$ .

So let us pick  $A \in S_{\varepsilon^2}$  and corresponding  $\vec{a} \in \mathbb{R}^3$  such that  $\|\vec{a}\| \leq \pi$  and  $A = e^{i\vec{a} \cdot \vec{\sigma}}$ . Now we choose  $\vec{b}, \vec{c} \in \mathbb{R}^3$  such that  $-2\vec{b} \times \vec{c} = \vec{a}$  and  $\|\vec{b}\| = \|\vec{c}\| = \sqrt{\|\vec{a}\|/2}$ . Define  $B = e^{i\vec{b} \cdot \vec{\sigma}}$  and  $C = e^{i\vec{c} \cdot \vec{\sigma}}$ .

Note that because  $\Gamma$  is an  $\varepsilon^2$ -net, that  $\|A - I\| = \|\vec{a}\| + O(\|\vec{a}\|^2) \leq \varepsilon^2$  and hence  $\|\vec{a}\| = O(\varepsilon^2)$ . By construction we then obtain  $\|\vec{b}\| = O(\varepsilon)$  and  $\|\vec{c}\| = O(\varepsilon)$ . We then invoke [Fact 4](#) to obtain

$$\|[B, C] - A\| = \left\| \left[ e^{i\vec{b} \cdot \vec{\sigma}}, e^{i\vec{c} \cdot \vec{\sigma}} \right] - e^{-[\vec{b} \cdot \vec{\sigma}, \vec{c} \cdot \vec{\sigma}]} \right\| = O(\varepsilon^3).$$

We now take  $U = e^{i\vec{u} \cdot \vec{\sigma}}, V = e^{i\vec{v} \cdot \vec{\sigma}} \in \Gamma$  to be the closest elements to  $B$  and  $C$  respectively, i.e.  $\|B - U\| \leq \varepsilon^2$  and  $\|C - V\| \leq \varepsilon^2$ . We can now use [Fact 1](#) to find

$$\|B - U\| = \|BU^\dagger - I\| = \left\| \vec{b} - \vec{u} \right\| + O\left( \left\| \vec{b} - \vec{u} \right\|^3 \right) \leq \varepsilon^2.$$

Therefore  $\|\vec{b} - \vec{u}\| = O(\varepsilon^2)$  and similarly  $\|\vec{c} - \vec{v}\| = O(\varepsilon^2)$ . We then obtain the following relation by the triangle inequality

$$\|A - \llbracket U, V \rrbracket\| \leq \left\| A - e^{2i(\vec{b} \times \vec{c}) \cdot \sigma} \right\| + \left\| e^{2i(\vec{b} \times \vec{c}) \cdot \sigma} - \llbracket U, V \rrbracket \right\|.$$

The right term is  $O(\varepsilon^3)$  by [Fact 4](#) so we consider the left term.

$$\begin{aligned} \left\| A - e^{2i(\vec{b} \times \vec{c}) \cdot \sigma} \right\| &\leq 2 \left\| \vec{b} \times \vec{c} - \vec{u} \times \vec{v} \right\| \\ &= 2 \left\| ((\vec{b} - \vec{u}) + \vec{u}) \times ((\vec{c} - \vec{v}) + \vec{v}) - \vec{u} \times \vec{v} \right\| \\ &= 2 \left\| (\vec{b} - \vec{u}) \times (\vec{c} - \vec{v}) + \vec{u} \times (\vec{c} - \vec{v}) + (\vec{b} - \vec{u}) \times \vec{v} \right\| \\ &= O(\varepsilon^4) + O(\varepsilon^3) + O(\varepsilon^3) \\ &= O(\varepsilon^3) \end{aligned}$$

Hence we have

$$\|A - \llbracket U, V \rrbracket\| = O(\varepsilon^3).$$

□

**Lemma 2.** Let  $\varepsilon > 0$  sufficiently small. There exists  $k \in \mathbb{R}$  such that if  $\Gamma$  is an  $\varepsilon^2$ -net for  $S_\varepsilon$  then  $\llbracket \Gamma, \Gamma \rrbracket \Gamma$  is a  $k^2\varepsilon^3$ -net for  $S_{k\varepsilon^{3/2}}$ .

*Proof.* By [Lemma 1](#) There exists  $k \in \mathbb{R}$  such that  $\llbracket \Gamma, \Gamma \rrbracket$  is a  $k^2\varepsilon^3$ -net for  $S_{\varepsilon^2}$ . Now pick  $A \in S_{k\varepsilon^{3/2}}$ . For sufficiently small  $\varepsilon$  we have  $S_{k\varepsilon^{3/2}} \subset S_{\varepsilon^2}$  so we know there exists  $W \in \Gamma$  such that  $\|AW^\dagger - I\| \leq \varepsilon^2$ . Thus we have  $AW^\dagger \in S_{\varepsilon^2}$  and hence we have  $U, V \in \Gamma$  such that  $\|AW^\dagger - \llbracket U, V \rrbracket\| \leq k^2\varepsilon^3$ .

□

Having established [Lemma 2](#) to reduce the size of our net, we apply the idea inductively. Establishing a means to create arbitrarily small nets around the identity. We formalise this in the following corollary.

**Corollary 1.** If  $\Gamma_0$  is an  $\varepsilon_0^2$ -net for  $S_{\varepsilon_0}$  sufficiently small. Then  $\Gamma_i = \llbracket \Gamma_{i-1}, \Gamma_{i-1} \rrbracket \Gamma_{i-1}$  is a  $\varepsilon_i^2$ -net where  $\varepsilon_i = (k^2\varepsilon_0)^{(3/2)^i} / k^2$  for some  $k \in \mathbb{R}$ .

Since each element of  $\Gamma_i$  is composed of gates in  $\Gamma_0$ , we effectively construct a sequence of gates; each taking us closer to the identity.

### 3 Solovay-Kitaev Theorem

**Theorem 1.** If  $\Gamma$  is a universal gate set that is closed under inverses. Then we can approximate any  $U \in SU(2)$  to any accuracy  $\varepsilon > 0$  by a sequence of gates in  $\Gamma$ .

*Proof.* Pick  $\varepsilon_0$  sufficiently small and independent of  $\Gamma$ . We wish to construct  $\Gamma_0$  - an  $\varepsilon_0^2$ -net for  $SU(2)$ . Since  $\langle \Gamma \rangle$  is dense, we can take the  $\varepsilon_0^2$  neighbourhoods of points in  $\langle \Gamma \rangle$  which form a cover for  $SU(2)$ . We now note that  $SU(2)$  is compact and hence has a finite subcover. Taking the centers of the subcover and their inverses we form  $\Gamma_0$ .

Pick  $U \in SU(2)$  and find  $V_0 \in \Gamma$  such that  $\|U - V_0\| = \left\| UV_0^\dagger - I \right\| \leq \varepsilon_0^2$ . We then have that  $UV_0^\dagger \in S_{\varepsilon_0^2}$ . For sufficiently small  $\varepsilon_0$  we have  $\varepsilon_0^2 < k\varepsilon_0^{3/2} = \varepsilon_1$  for the  $k \in \mathbb{R}$  given by [Lemma 1](#). Thus by [Corollary 1](#), we have that  $\Gamma_1$  is an  $\varepsilon_1^2$ -net for  $S_{\varepsilon_1}$  and we can find  $V_1 \in \Gamma_1$  such that

$$\left\| UV_0^\dagger V_1^\dagger - I \right\| = \|U - V_1 V_0\| \leq \varepsilon_1^2 < k\varepsilon_1^{3/2} = \varepsilon_2.$$

Proceeding inductively, we can find  $V_t \in \Gamma_t$  with  $\|U - V_t \cdots V_0\| \leq \varepsilon_t^2$ . We note that each  $V_i$  is composed of 5<sup>i</sup> gates. And hence we need  $\sum_{i=0}^t 5^i = O(5^t)$  gates in  $\Gamma_0$ . Moreover, for an accuracy  $\varepsilon$  we need  $\varepsilon_t^2 = ((k^2\varepsilon_0)^{(3/2)^t}/k^2)^2 \leq \varepsilon$  and solve for  $t$ . Seeing as  $\frac{3}{2}^{(\log(5)/\log(3/2))} = 5$ . We set  $c = \frac{\log(5)}{\log(3/2)}$  and we find,

$$\begin{aligned} ((k^2\varepsilon_0)^{(3/2)^t})^2 &\leq k^4\varepsilon \\ \left(\frac{3}{2}\right)^t \ln(k^2\varepsilon_0) &\geq \frac{1}{2} \ln(k^4\varepsilon) \\ \left(\frac{3}{2}\right)^t &\leq \frac{\log(k^4\varepsilon)}{2 \log(k^2\varepsilon_0)} \\ \left(\frac{3}{2}\right)^t &\leq \frac{\log(1/k^4\varepsilon)}{2 \log(1/k^2\varepsilon_0)} \\ 5^t &\leq \frac{\log^c(1/k^4\varepsilon)}{2 \log^c(1/k^2\varepsilon_0)} \end{aligned}$$

Hence we have  $O(5^t) = O(\log^c(1/\varepsilon))$ . □

## 4 Results in Higher Dimensions

As our results hold only for gates in  $SU(2)$ , and therefore for single qubit systems, we wish to extend the proof for arbitrary qudits. I.e, a proof that holds for  $SU(d)$ .

Such a result is given in a paper by Dawson and Nielsen [\[2\]](#). Where an  $\varepsilon$ -approximate can be made with the paradoxically better result of  $O(\log^{2.71}(1/\varepsilon))$  gates because of some redundancy that occurs. If one chooses to ignore this redundancy, a similar proof to the above can be written, giving the familiar bound of  $O(\log^{3.97}(1/\varepsilon))$ .

### 4.1 Facts about $SU(d)$

The proof given in this report generalises to  $SU(d)$  to the extent that the identities given change slightly but the idea of the proof is the same. For example [Fact 1](#) in  $SU(d)$  becomes:

**Fact 5.** If  $H$  is hermitian then

$$\|e^{iH} - I\| = 2 \sin\left(\frac{\|H\|}{2}\right) = \|H\| + O(\|H\|^3).$$

## 5 Contains an IRREP

In our definition of universal gate set, we insisted that it be closed under inverses. A paper by Bouland and Ozols [3] proves that one can instead insist that our set  $\Gamma$  is instead contains a projective irreducible representation of some finite group  $G$ . To approximate any  $U \in SU(d)$  to accuracy  $\varepsilon$ , one only needs  $O(\log^{\log_2 |G|}(1/\varepsilon))$  gates in  $\Gamma$ .

### 5.1 Representations of Groups

An representation of a group  $G$  over vector space  $\mathbb{C}^n$  is a homomorphism  $\sigma : G \rightarrow GL(\mathbb{C}^n)$ . Given two representations  $\sigma : G \rightarrow V$  and  $\sigma' : G \rightarrow V'$  we define their direct sum  $\sigma \oplus \sigma' : G \rightarrow V \oplus V'$  by

$$(\sigma \oplus \sigma')(x) = \begin{pmatrix} \sigma(x) & 0 \\ 0 & \sigma'(x) \end{pmatrix}$$

for all  $x \in G$ . A representation is then irreducible if it is not the direct sum of two other representations. More specifically, this paper makes use of representations  $\sigma : G \rightarrow U(d)$  and calls  $\sigma$  projective if there is some function  $\theta : G \times G \rightarrow \mathbb{R}$  such that for all  $g_1, g_2 \in G$   $\sigma(g_1)\sigma(g_2) = e^{i\theta(g_1, g_2)}\sigma(g_1, g_2)$ . i.e,  $\sigma$  preserves global phase.

The theorem then becomes

**Theorem 2.** For any fixed  $d \geq 2$ , suppose  $\Gamma \subset SU(d)$  is a finite gate set which densely generates  $SU(d)$ , and furthermore  $\Gamma$  contains a (projective) irrep of some finite group  $G$ . Then there is an algorithm which outputs an  $\varepsilon$ -approximation to any  $U \in SU(d)$  using merely  $O(\text{polylog}(1/\varepsilon))$  elements from  $\Gamma$ .

## References

- [1] Maris Ozols. The Solovay-Kitaev theorem. page 9.
- [2] Christopher M. Dawson and Michael A. Nielsen. The Solovay-Kitaev algorithm. *arXiv:quant-ph/0505030*, May 2005. arXiv: quant-ph/0505030.
- [3] Adam Bouland and Maris Ozols. Trading inverses for an irrep in the Solovay-Kitaev theorem. *arXiv:1712.09798 [math-ph, physics:quant-ph]*, page 15 pages, 2018. arXiv: 1712.09798.